

Stacks Seminar - Scribe Notes for Jeremy N's Talk 3/21/25

Descent Theory - Faithfully Flat Descent and the Stack Condition

n.b. we will work with the affine site $(\text{Aff}/S)_{\text{fppf}}$
 it gives the same shv (see Olsson; Stacks §34.7), to $(\text{Sch}/S)_{\text{fppf}}$.

Motivation: Gluing Lemma for Sheaves on Top Space X .

Recall def of Sheaf on X : a contravariant functor $\text{Open}(X) \xrightarrow{F} \text{Set}$
 satisfying two conditions

i) (Separation) For $\coprod_{i \in I} U_i \rightarrow U$ open cover, $s, t \in F(U)$
 and s_i, t_i be image of s, t in $F(U_i)$.

If $s_i = t_i$ $\forall i$, then $s = t$. (called sep presheaf)

ii) (gluing) For such coverings, $s_i \in F(U_i)$ we have

$$s_{ij} \triangleq s_i|_{U_i \cap U_j}, \quad s_{ij} = s_{ji}$$

then $\exists! s \in F(U)$ s.t. $s|_{U_i} = s_i$.

Gluing Lemma: If $\coprod U_i$ covering of X , F_i sheaves on U_i , s.t.

- $\varphi_{ij} : F_i|_{U_i \cap U_j} \xrightarrow{\sim} F_j|_{U_i \cap U_j}$

- and $\varphi_{ik} = \varphi_{jk} \circ \varphi_{ij}$ on $U_i \cap U_j \cap U_k$,

then $\exists!$ F on X sheaf s.t. $F|_{U_i} = F_i$.

this
must
change in F.C. setting.

Remark: We want to generalize to fibred categories on a site, in particular, to do this to QCoh. "gluing" is what we call "descent".

Descent for modules

Lemma: modules form a sheaf on $(\text{Aff}/\text{Spec } A)_{\text{fppf}}$.

$F: A \rightarrow B$ is fppf, M an A -module, then

$$(*) \quad 0 \rightarrow M \rightarrow M \otimes_A B \xrightarrow{p_1 - p_2} M \otimes_A B \otimes_B B \quad \text{is exact.}$$

$m \mapsto m \otimes 1$

$A\text{-modules}$

(Reminder: faithfully flat means $0 \rightarrow M_1 \rightarrow M_2 \rightarrow M_3 \rightarrow 0$ exact $A\text{-mod}$)
 $\Leftrightarrow 0 \rightarrow M_1 \otimes_B B \rightarrow M_2 \otimes_B B \rightarrow M_3 \otimes_B B \rightarrow 0$

$$\text{Here, } p_1(m \otimes b) = m \otimes b \otimes 1$$

$$p_2(m \otimes b) = m \otimes 1 \otimes b$$

(Grothendieck trick)

Proof: Suppose a section g of F exists i.e. $g \circ F = \text{id}_A$.

Then $(*)$ exact at M i.e. $M \hookrightarrow M \otimes_A B$.

Exactness at $M \otimes_B B$. Say $m \otimes b \in \ker(p_1 - p_2)$. Then

$$m \otimes b \otimes 1 - m \otimes 1 \otimes b = 0$$

$$\Leftrightarrow m \otimes (b \otimes 1 - 1 \otimes b) = 0$$

$$\Leftrightarrow b \otimes 1 - 1 \otimes b = 0$$

$$\Rightarrow b \in \text{Im of } A.$$

* exercise

$$\xrightarrow{\text{b/c}} 0 \rightarrow A \rightarrow B \rightarrow B \otimes_B B \quad \text{is exact.}$$

□

Corollary: (Descent for modules) $F: A \rightarrow B$ ff.

$$M' - B\text{-module}, \quad \psi: M' \otimes_A B \xrightarrow{\sim} B \otimes_A M'$$

$$B \otimes_A M' \otimes_B B \xrightarrow{\sim} B \otimes_B B \otimes_A M' \xleftarrow{\sim} M' \otimes_B B \otimes_B B \xrightarrow{\sim} B \otimes_B M' \otimes_B B$$

Suppose $\varphi_1 \circ \varphi_3 = \varphi_2$. Then $\exists! A\text{-mod } M' \text{ s.t. } M \otimes B = M'$.

Pictorially:

$$\begin{array}{ccc} M' & \xrightarrow{\quad} & M \\ \downarrow \varphi_1 \circ \varphi_3 & \xrightarrow{\quad} & \downarrow \varphi_2 \\ M & & \end{array}$$

$\text{Spec } B \otimes_A B \rightarrow \text{Spec } B \rightarrow \text{Spec } A$

" $\text{Spec } B \cap \text{Spec } B'$ "

or just (M, φ)

Say $(M', \varphi_1, \varphi_2, \varphi_3)$ the descent data.

$$\begin{array}{ll} \alpha: M \rightarrow B \otimes M' & m \mapsto 1 \otimes m \\ \beta: M \rightarrow B \otimes M' & m \mapsto \varphi(m \otimes 1) \end{array}$$

Proof:

$$M' \otimes B \rightarrow B \otimes M' \otimes B$$

$$\begin{array}{ccc} \downarrow \varphi & & \downarrow \varphi_1 \\ B \otimes M' & \xrightarrow{\quad} & B \otimes B \otimes M' \\ p_1 - p_2 & & \end{array}$$

By $\varphi_1 \circ \varphi_3 = \varphi_2$, the diagram commutes.

By faithful flatness, kernel of $\alpha \otimes 1 - \beta \otimes 1 = M \otimes_A B$ where

$$M = \{m' \in M \mid 1 \otimes m - \varphi(m \otimes 1) = 0\}.$$

Kernel of $p_1 - p_2 \rightarrow M'$. 5-lemma says $M \otimes_A B \cong M'$.

Theorem: Suppose $F: \text{Spec } B \rightarrow S$ is a covering in $(\text{Affts})_{\text{fppf}}$.

Let $(M', \varphi_1, \varphi_2, \varphi_3)$ Descent data for \tilde{M}' on $\text{Spec } B$.

Then $\exists!$ qcoh sheaf N on S s.t. $F^*N = \tilde{M}'$.

Corollary: $(\mathcal{Q}\text{Coh})^{\text{pf}}, \text{Coh}, \text{Vect}$ all have fppf descent.

Affine morphisms, closed subschemes satisfy descent.

Corollary: Polarized schemes have descent.

Pol obs $(f: X \rightarrow U, L)$

\downarrow U affine L f -relatively ample.

$(\text{Aff}/S)_{\text{fppf}}$

f proper.

Mors $(f, L) \rightarrow (f', L')$

are

someone asked what this - excellence should look like.

$f' \downarrow \begin{matrix} X' \xrightarrow{u} X \\ f' \downarrow \lrcorner \\ U' \xrightarrow{u'} U \end{matrix}$

$X' \xrightarrow{u} X$

$f' \downarrow \lrcorner$

$U' \xrightarrow{u'} U$

+ an iso $u^* L \xrightarrow{\sim} L'$.

It's fibred category $(\text{Aff}/S)_{\text{fppf}}$

by

$(f, L) \longmapsto U$.

Suppose $U' \rightarrow U$ an fppf covering in $(\text{Aff}/S)_{\text{fppf}}$.

By proper flat base change $u'^* f'_* L' \cong f'_* u^* L \cong f'_* L'^{\otimes N}$

$$\Rightarrow u'^* E = E'$$

\Rightarrow projective embedding base changes i.e.

$$X' \hookrightarrow \mathbb{P}(E')$$

$$X \hookrightarrow \mathbb{P}(E)$$

descent follows from for loc free \neq closed subschemes.

Corollary: M_g , $g \geq 1$ and A_g satisfy descent (objects of M_g are pol curves pol by w , $g \geq 2$)
 (objects of A_g are polarized by Θ -divisor). $w^{-1}, g=0$)

Def: $F \xrightarrow{p} C$ be a cat fibred in groupoids, C site.

$x, y \in F(U)$, define $\text{Isom}(x, y) : C/U \rightarrow \text{Set}$
 $V \xrightarrow{f} U \longmapsto \text{Isom}(f^*x, f^*y).$

this is a presheaf on C/U .

Def: A stack in groupoids is a CFG^v s.t.

$F \xrightarrow{(\text{AffS})_{\text{perf}}}$

(1) $\text{Isom}(x, y)$ are sheaves $\forall x, y \in F(U), \forall U \in (\text{AffS})_{\text{perf}}$.

(2) Every descent data is effective.

i') Isomorphisms are separated

ii') Isomorphisms glue.

Expl: $F = \begin{matrix} \text{cat of} \\ \text{isomorphic} \\ \text{line bundles} \end{matrix}$ i.e. $U \mapsto \text{Pic}(U)$. This is not separated.

Expl: Most examples from before can be made into stacks.

1) $\mathbb{Q}\text{Coh}$, $\mathbb{Q}\text{Coh}^{pf}$, and Coh are all stacks.
 i.e. $U \mapsto (-)(U)$.

2) $Mg \neq Ag$ ($g \neq 1$ for Mg).

3) Bun_G , and relative vector bundles. Vect_X .

left adjoint to
forgetful

Stackification: There's a functor $(-)^a : \text{CFG} \rightarrow \text{Stacks}$
 satisfying universal property $\underset{\text{CFG}}{\text{HOM}}(F, G) \xrightarrow{\sim} \underset{\text{Stacks}}{\text{HOM}}(F^a, G)$.
 G a stack

Expl: quotient stack let X be a scheme, w/ G -action, \underline{G} smooth affine group scheme

$[X/G]$ objects over $U = \text{Hom}_S(U, X)$

maybe not
need.

morphisms = $\{g \in G \mid g \circ u = u'\}$.

The particular case $X = *$ is called the classifying stack BG of G .